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## **Optimal Closing Benchmarks**

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# Optimal Closing Benchmarks

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## Abstract

In financial markets, the closing price serves as an important benchmark. Accordingly, its stability is of great concern to market administrators, as they desire to have a benchmark that is robust against distorting and manipulating trades. We introduce a market model to analyze the stability of the closing price with presence of three types of volume: distorting volume, volume that targets the closing price, and volume that is unrelated to the closing price. We find that the optimal closing price is either the price from an auction or the volume weighted average price (VWAP) from regular trading only, explaining the prevalence of these closing benchmarks on financial markets. A succinct condition depending on the different volume types indicates when the inclusion of a closing auction is optimal.

**Keywords:** benchmark stability, closing auction, volume weighted average price, VWAP

**JEL Codes:** G14, G18, D44, D82

## 1 Introduction

Closing prices are important benchmarks in financial markets since many investors use them as reference points. In practice, two methods to determine the closing price are predominant: a closing auction, used for stocks at all developed markets and some emerging markets, and the volume weighted average price (VWAP), common in futures markets and several emerging stock markets; see FTSE Russell [2019]. A primary concern with benchmark design is robustness to manipulating or distorting orders. For example, the main reason for the Paris

Bourse (currently Euronext Paris) to move to a closing auction was that some relatively small orders could have changed closing prices, as Pagano and Schwartz [2003] note. The stability issue of closing prices is underlined by judicial and empirical evidence for their manipulation [Commodity Futures Trading Commission, 2013, Hillion and Suominen, 2004].

The closing price of a stock is especially important to mutual funds as purchases and redemptions of their shares occur at the day's closing price of the funds' constituent assets. This creates the desire for mutual funds to buy and sell stocks at their closing prices. As a result, there is a significant portion of volume that follows prices used to create the benchmarks. This is especially noticeable in markets after the introduction of a closing auction.<sup>1</sup> Kandel et al. [2012] as well as Hagströmer and Nordén [2014] find that trader patience increases after the introduction of a closing auction as a result of improved liquidity during the end of regular trading. It would then be prudent to consider the effect of a benchmark change that results from such benchmark targeting volume.

We include this volume as a feature of our market model, which also has linear price impact during regular trading and distorting volume. We take the perspective of a benchmark administrator, who wishes to minimize the impact of distorting volume on the benchmark. The administrator decides (1) whether or not to include an auction in the market, (2) the benchmark weight of the auction if it is included, and (3) the distribution of the remaining weights in the regular trading. We find a unique optimal benchmark that minimizes the impact of distorting volume under worst-case scenario assumptions, as well as a succinct condition for when the inclusion of a closing auction in this market is optimal. We also find that in the exclusion of a closing auction, the unique optimal benchmark is VWAP and give a method to find the optimal start time of the window, over which VWAP is computed.

On the one hand, our result explains the prevalence of VWAP and auction benchmarks in practice; on the other hand, it gives a criterion when VWAP or an auction is preferable, even determining the optimal window length in the case of VWAP. This condition can be applied to a market with an estimation of benchmark targeting volume relative to volume that is unrelated to the benchmark and provides insight into why these two kinds of benchmarks appear in different markets.

VWAP-type benchmarks have been found to be optimal in different settings and questions. For example, Duffie and Dworczak [2018] examine optimal benchmarks in settings where data from transactions or reports of agents whose profits depend on such benchmarks are available. They find a benchmark that puts small weight on small transactions and nearly equal weight on all large transactions. When such data is unavailable, VWAP emerges as

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<sup>1</sup>For large developed markets of stocks, closing auctions account for as much as 20% of the daily traded volume [Golden, 2018].

the optimal benchmark. Baldauf et al. [2019] find VWAP as the unique optimal benchmark for the principal agent problem of a client contracting a broker to purchase a large amount of shares. While a closing auction based benchmark may seem more attractive to agents in the marketplace who target the benchmark, execution algorithms with VWAP benchmarks are readily available and well studied [Cartea and Jaimungal, 2016, Frei and Westray, 2015, Guéant and Royer, 2014, Humphery-Jenner, 2011, Kato, 2015].

A VWAP benchmark includes volume that is unrelated to the specific benchmark, providing additional stability. However, closing auctions can help absorb liquidity shocks by pooling volume. Foucault et al. [2005] and Roşu [2009] study dynamic models of limit order markets with strategic liquidity traders of varying patience. Their conclusions indicate that the introduction of a closing auction can increase trader patience and thus reduce spreads. These theoretical findings are confirmed by Pagano and Schwartz [2003] and Kandel et al. [2012] for the stock markets of the Borsa Italiana and the Paris Bourse, where they document a significant reduction in spread and volatility around the close when closing auctions were introduced. For NASDAQ-OMXS30 index futures, which have a closing auction as an exception among futures markets, Hagströmer and Nordén [2014] find that the introduction of a closing auction improved closing price accuracy.

The remainder of this paper is organized as follows. Section 2 describes a market model with linear price impact during regular trading and three types of volume: benchmark targeting volume, distorting volume, and noise volume that is unrelated to the benchmark. We use the hypothetical shape of an order book under linear price impact to build supply and demand curves present in an auction and derive the price impact of an order submitted to the auction. Section 3 introduces the optimization that the benchmark administrator faces and states the main result, that is, the unique optimal benchmark and the aforementioned condition for auction inclusion. Section 4 contains the proof of the main result.

## 2 The market model

We consider a discrete market model with  $T$  periods consisting of  $T - 1$  periods of regular trading and possibly an auction at time  $T$ .<sup>2</sup> We will later introduce a benchmark administrator, who decides about the closing benchmark and whether or not an auction takes place. The market is made up of three types of volume:

- Noise or “outside” volume, denoted as  $u_i$ , which is unsigned ( $u_i \geq 0$ ) and does not follow the benchmark.

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<sup>2</sup>Alternatively, we could consider a continuous-time model for regular trading. Our results continue to hold in such a continuous-time version under a specification analogous to the discrete model that we present.

- Benchmark targeting volume, denoted as  $\alpha_i$  and summing to  $A$ ; this is the result of traders targeting the benchmark price. Here,  $\alpha_i$  and  $A$  also denote unsigned volumes so that  $\alpha_i \geq 0$  and  $A \geq 0$ . We assume that traders responsible for this volume can essentially perfectly match the benchmark and let  $\alpha_i = A\beta_i$ .<sup>3</sup>
- Distorting volume, denoted as  $v_i$  and summing to  $V$ ; this volume is made up of potential distortions to the benchmark. We assume that all  $v_i$  have the same sign.

All of  $u_i$ ,  $\alpha_i$ ,  $v_i$ ,  $A$ , and  $V$  are random variables. The prices observed are modelled as

$$p_i = \tilde{p}_i + c \frac{v_i}{u_i + \alpha_i} \quad \text{for } i = 1, \dots, T - 1.$$

Such a model of temporary price impact is well founded in the literature. Indeed, linear price impact (in  $v_i$ ) is consistent with price impact models based on adverse selection [Kyle et al., 2018]. The ratio  $\frac{v_i}{u_i + \alpha_i}$  makes the price impact dimensionless [Almgren et al., 2005]. The linear price impact model in regular trading can be seen as each period having an order book with a constant amount of sell orders above  $\tilde{p}_i$  and a symmetric amount of buy orders below of  $\tilde{p}_i$ . Indeed, the price changes by one unit for each volume of  $\frac{u_i + \alpha_i}{c}$ . A continuous version of this order book is shown below in Figure 1.

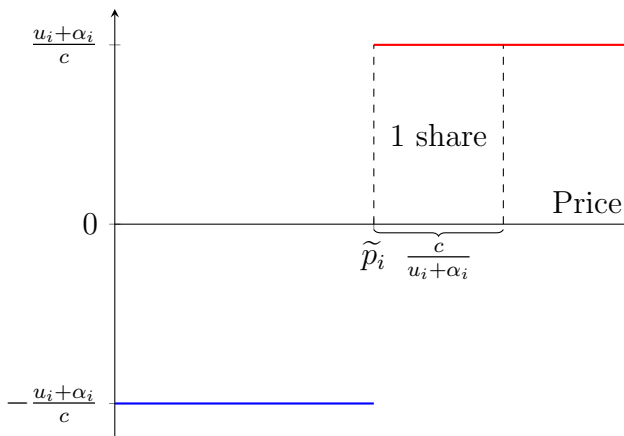


Figure 1: the order book such that buying  $v_i$  shares increases the price by  $c \frac{v_i}{u_i + \alpha_i}$

In the auction, when all volume is pooled, we can represent the price as the intersection of supply and demand curves. We assume that volume unrelated to the auction,  $u_T$ , is zero. The supply and demand curves correspond to the aggregation of the sell and buy orders, respectively. Because our model is based on constant amounts of sell and buy orders at

<sup>3</sup>In the optimum, the benchmark will be either the auction price or the VWAP, depending on the relative volume curve, which can be forecast to a fair degree of accuracy (e.g., Exhibit 9 of Satish et al. [2014]).

different prices, this aggregation leads to supply and demand curves that are linear and have the same slope in absolute value, as the integration of a constant results in a linear function. A distorting buy order will shift the demand curve to the right while a distorting sell order will shift the supply curve to the right; compare Figure 2.

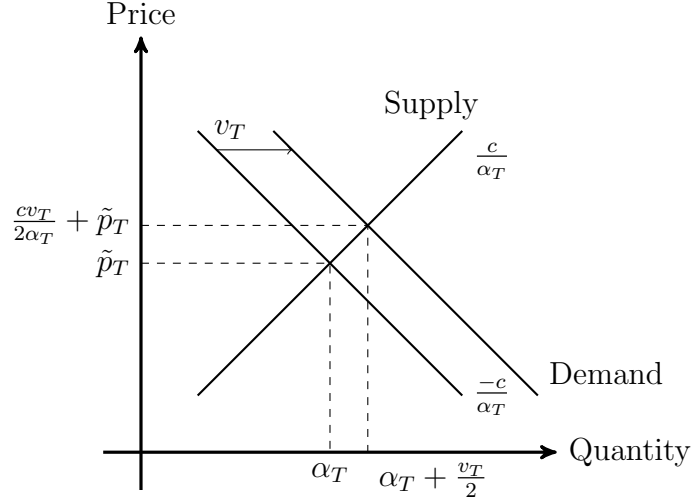


Figure 2: the shifting of the demand curve in the auction from a distorting buy order  $v_T$

Since the slopes of the supply and demand lines are the negative of one another, shifting the demand curve by  $v_T$  causes the new equilibrium quantity to shift by only  $v_T/2$ , which results in a price impact of  $c\frac{v_T}{2\alpha_T}$ . Consequently, we model the auction price as

$$p_T = \tilde{p}_T + c\frac{v_T}{2\alpha_T},$$

provided that the benchmark administrator decides to have an auction.

### 3 The benchmark administrator's optimization

The administrator desires to choose a benchmark such that the impact of the distorting orders,  $v_i$ , on the benchmark is minimized. The administrator can choose a nonnegative weight  $\beta_i$  for trading period  $i$ , where we normalize  $\sum_{i=1}^T \beta_i = 1$ . Since an auction consists only of volume related to the benchmark, choosing  $\beta_T = 0$  is equivalent to having no auction. The impact of the distorting orders on the benchmark is given by  $\sum_{i=M}^T \beta_i(p_i - \tilde{p}_i)$ . Because the benchmark is used at the close, the administrator is also worried about the cost that comes from starting early with nonzero benchmark weights. This cost can be related to the need of a longer time interval for computing the benchmark, which leads to additional complexity and means clients of mutual funds would not be able to submit orders during

this time. Let  $M$  be the first nonzero weight of the benchmark so that  $\beta_M \neq 0$  and  $\beta_i = 0$  for all  $i < M$ . We model the costs of starting the benchmark computation early by  $Q(M)$ , where  $Q$  is a strictly decreasing function.

To see why it is desirable to have the benchmark start later in the day, consider a mutual fund whose inflows occur at the benchmark and a benchmark that starts in the last 10 periods of trading. If an underlying asset in the mutual fund's price takes a significant jump with 5 periods left in the day, one could purchase the mutual fund and obtain the constituent asset at a price that is based on the last 10 periods and likely to be less than the closing price. To combat this, mutual funds typically do not allow trades during the benchmark building period.

Thus, the benchmark administrator faces the optimization

$$\min_{\beta_i, M} \left\{ \max_{v_i} E \left[ \left| \sum_{i=M}^T \beta_i (p_i - \tilde{p}_i) \right| + Q(M) \right] \right\}, \quad (1)$$

where the maximization is over all random variables  $v_i$  such that  $\sum_{i=1}^T v_i = V$  with all  $v_i$  having the same sign as  $V$ , and the minimization is over  $M$  and all  $\beta_i$  that are functions of the prices and volumes for given  $M$  and satisfy  $\sum_{i=M}^T \beta_i = 1$  with  $\beta_i = 0$  for  $i < M$ . Mathematically, the condition that  $\beta_i$  are functions of the prices and volumes for given  $M$  means that  $\beta_i$  are measurable with respect to the  $\sigma$ -algebra generated by  $p_M, \dots, p_{T-1}$  and  $u_M + \alpha_M, \dots, u_{T-1} + \alpha_{T-1}$ . As  $\alpha_i = A\beta_i$  themselves depend on  $\beta_i$ , this circular condition may seem mathematically complicated, but we will show that the optimal benchmark weights are well-defined and satisfy this measurability condition.<sup>4</sup>

We denote by  $M^*$  be the minimizer of

$$E \left[ \frac{c|V|}{A + \max\{\sum_{j=M}^{T-1} u_j, A\}} + Q(M) \right] \quad (2)$$

over  $\{1, 2, \dots, T\}$ .

**Main Result.** *There exists a unique optimal benchmark.*

*If  $M^* = T$ , there is a closing auction and the optimal benchmark puts all weight into the closing auction with  $\beta_T^* = 1$  and  $\beta_i^* = 0$  for all  $i < T$ .*

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<sup>4</sup>If  $A$  is deterministic the condition is equivalent to requiring that  $\beta_i$  are measurable with respect to the  $\sigma$ -algebra generated by  $p_M, \dots, p_{T-1}$  and  $u_M, \dots, u_{T-1}$ . For stochastic  $A$ , however, the benchmark administrator may not be able to distinguish between the noisy volume  $u_i$  and the volume  $\alpha_i$  that follows the benchmark.

If  $M^* < T$ , there is no closing auction and the optimal benchmark is VWAP with  $\beta_j^* = \frac{A\beta_j^* + u_j}{A + \sum_{\ell=M^*}^{T-1} u_\ell} = \frac{u_j}{\sum_{\ell=M^*}^{T-1} u_\ell}$  for  $j = M^*, \dots, T-1$  and  $\beta_T^* = 0$  and  $\beta_i^* = 0$  for  $i < M^*$ .

We see that the decision of whether or not to have an auction is a simple yes or no. This is consistent with what we see in practice: VWAP and auction closing prices are predominant, but no market computes the closing price as a weighted combination of VWAP and the auction price. The decision between VWAP and closing auction as well as when to start the VWAP window, can be tweaked depending on a benchmark administrator's perceived penalty for starting the benchmark period early through the choice of  $Q(M)$ .

In practice, our main result can be applied in two steps. First,  $M^*$  will be estimated based on patterns observed in historical volume. This will then lead to the choice of either a closing auction or a VWAP benchmark. In the latter case, the weights of the benchmark will be computed each day after the close based on the observed volume. For this second step, it is useful that the weights depend on the total volume  $u_i + \alpha_i$  in each period  $i$ , but not on the split of this volume between  $u_i$  and  $\alpha_i$ .

In the case  $\sum_{j=1}^{T-1} u_j \leq A$  almost surely, that is less volume that is unrelated to the benchmark than there is targeting it, we can see that the ratio  $\frac{c|V|}{A + \max\{\sum_{j=M}^{T-1} u_j, A\}}$  in (2) does not depend on  $M$ . Because  $Q$  is strictly decreasing, it is then optimal to choose  $M^* = T$ , and we can deduce from our main result the following corollary.

**Corollary.** *If  $\sum_{j=1}^{T-1} u_j \leq A$  almost surely, there is a closing auction and the optimal benchmark puts all weight into the closing auction with  $\beta_T^* = 1$  and  $\beta_i^* = 0$  for all  $i < T$ .*

By contrast, if the distributions of  $\sum_{j=1}^{T-1} u_j$  and  $A$  overlap or even if  $\sum_{j=1}^{T-1} u_j > A$  almost surely, then one must choose  $Q(M)$  explicitly and optimize (2) over  $M^*$  either analytically or numerically to determine if a closing auction is optimal or not. If we have that  $\sum_{j=1}^{T-1} u_j > A$  almost surely and  $Q$  has a small slope in absolute value such that its marginal impact on (2) is minimal, then VWAP will be the optimal benchmark.

Generally, the bigger the benchmark targeting volume is relative to the volume that is unrelated to the closing price, the more likely the minimizer of (2) will be  $T$  so that a closing auction is optimal. Interestingly, we observe lower volume around the close in emerging stock markets compared to developed stock markets, which saw a rise in volume around the close [Golden, 2018]. This observation motivates why developed stock markets moved to closing auctions over the last two decades while many emerging stock markets still use a VWAP benchmark. Our main result gives in (2) a concrete criterion when a VWAP benchmark or a closing auction are optimal.



## 4 Proof of the main result

*Step 1: inner maximization problem*

We first write the inner maximization problem in (1) as

$$\max_{v_i} E \left[ \left| \sum_{i=M}^T \beta_i (p_i - \tilde{p}_i) \right| \right] = c \max_{v_i} E \left[ \left| \sum_{i=M}^{T-1} \beta_i \frac{v_i}{u_i + \alpha_i} + \beta_T \frac{v_T}{2\alpha_T} \right| \right]. \quad (3)$$

We will use the auxiliary result that for any nonnegative  $\kappa_M, \dots, \kappa_T$ , we have

$$\max_{v_i} \left| \sum_{i=M}^T \kappa_i v_i \right| = |V| \max_{j=M, \dots, T} \kappa_j. \quad (4)$$

This result follows from

$$\left| \sum_{i=M}^T \kappa_i v_i \right| \leq \left| \max_{j=M, \dots, T} \kappa_j \sum_{i=M}^T v_i \right| = |V| \max_{j=M, \dots, T} \kappa_j,$$

where we used for the inequality that all  $v_i$  have the same sign. Equality holds if  $v_{i^*} = V$  for  $i^* = \operatorname{argmax}_{i=M, \dots, T} \kappa_i$  (if there are several  $i$  with maximal  $\kappa_i$ , we can choose  $i^*$  arbitrarily among them) and  $v_i = 0$  for all  $i \neq i^*$ .

Using that the worst case of an expectation is given by the worst case in each scenario, we apply (4) to (3) in each scenario with  $\kappa_i = \frac{\beta_i}{u_i + \alpha_i}$  for  $i = M, \dots, T-1$  and  $\kappa_T = \frac{\beta_T}{2\alpha_T}$ , which yields

$$\max_{v_i} E \left[ \left| \sum_{i=M}^T \beta_i (p_i - \tilde{p}_i) \right| \right] = c E \left[ |V| \max \left\{ \max_{j=M, \dots, T-1} \frac{\beta_j}{u_j + \alpha_j}, \frac{\beta_T}{2\alpha_T} \right\} \right]. \quad (5)$$

*Step 2: outer minimization problem*

Thanks to (5) and  $\alpha_i = A\beta_i$ , the optimization problem (1) becomes

$$\min_{\beta_i, M} \left\{ E \left[ c |V| \max \left\{ \max_{j=M, \dots, T-1} \frac{\beta_j}{u_j + A\beta_j}, \frac{1}{2A} \mathbb{1}_{\beta_T \neq 0} \right\} + Q(M) \right] \right\}. \quad (6)$$

*Step 2a: outer minimization problem over  $\beta_M, \dots, \beta_{T-1}$*

For given  $M < T$  and  $\beta_T < 1$ , we first analyze the minimization of

$$\max_{j=M, \dots, T-1} \frac{\beta_j}{u_j + A\beta_j}$$

over  $\beta_M, \dots, \beta_{T-1}$ . We note that  $\frac{\beta_j}{u_j + A\beta_j}$  is increasing in  $\beta_j$  since differentiating with respect

to  $\beta_j$  shows that

$$\frac{\partial}{\partial \beta_j} \frac{\beta_j}{u_j + A\beta_j} = \frac{u_j}{(u_j + A\beta_j)^2} > 0.$$

In the minimum over  $\beta_M, \dots, \beta_{T-1}$ , we need equality of the ratios

$$\frac{\beta_i}{u_i + A\beta_i} = \frac{\beta_j}{u_j + A\beta_j} \quad \text{for all } i, j = M, \dots, T-1,$$

as otherwise, we could make  $\max_j \frac{\beta_j}{u_j + A\beta_j}$  smaller by making some  $\beta_j$  smaller and adding weight to another  $\beta_i$ . This leads to the optimal  $\beta_j^*$  of the form

$$\beta_j^* = a(u_j + A\beta_j^*),$$

where  $a$  is a random variable not depending on  $j$ . Using the constraint that  $\sum_{j=M}^{T-1} \beta_j^* = 1 - \beta_T$ , we deduce

$$\sum_{j=M}^{T-1} \beta_j^* = \sum_{j=M}^{T-1} a(u_j + A\beta_j^*) = 1 - \beta_T,$$

which gives the solution

$$\beta_j^* = \frac{u_j + A\beta_j^*}{A + \sum_{i=M}^{T-1} u_i / (1 - \beta_T)} \quad (7)$$

since

$$\sum_{i=M}^{T-1} A\beta_i^* = A(1 - \beta_T).$$

The above definition of  $\beta_j^*$  is equivalent to

$$\beta_j^* = \frac{u_j}{\sum_{i=M}^{T-1} u_i} (1 - \beta_T), \quad (8)$$

which shows that the optimal  $\beta_j^*$  for  $j = M, \dots, T-1$  are unique when  $M < T$  and  $\beta_T < 1$ . Moreover, we see from (7) and (8) that when there is no closing auction, which means  $\beta_T = 0$ , then the optimal benchmark is VWAP with  $\beta_j^* = \frac{A\beta_j^* + u_j}{A + \sum_{i=M}^{T-1} u_i} = \frac{u_j}{\sum_{i=M}^{T-1} u_i}$ .

*Step 2b: outer minimization problem over  $\beta_T$  and  $M$*

Thanks to (7), the optimization problem (6) becomes

$$\min_{\beta_T, M} \left\{ E \left[ c|V| \max \left\{ \frac{1}{A + \sum_{i=M}^{T-1} u_i / (1 - \beta_T)} \mathbb{1}_{\beta_T \neq 1}, \frac{1}{2A} \mathbb{1}_{\beta_T \neq 0} \right\} + Q(M) \right] \right\}. \quad (9)$$

We can write

$$\begin{aligned}
& \max \left\{ \frac{1}{A + \sum_{i=M}^{T-1} u_i / (1 - \beta_T)} \mathbb{1}_{\beta_T \neq 1}, \frac{1}{2A} \mathbb{1}_{\beta_T \neq 0} \right\} \\
&= \frac{1}{A} \max \left\{ \frac{1 - \beta_T}{1 - \beta_T + \sum_{i=M}^{T-1} u_i / A} \mathbb{1}_{\beta_T \neq 1}, \frac{1}{2} \mathbb{1}_{\beta_T \neq 0} \right\} \\
&= \frac{1}{A} \left( \underbrace{\frac{1}{1 + \sum_{i=M}^{T-1} u_i / A}}_{\text{term 1}} \mathbb{1}_{\beta_T=0} + \mathbb{1}_{\beta_T \neq 0} \underbrace{\max \left\{ \frac{1 - \beta_T}{1 - \beta_T + \sum_{i=M}^{T-1} u_i / A}, \frac{1}{2} \right\}}_{\text{term 2}} \right). \tag{10}
\end{aligned}$$

We analyze the following two cases separately:

1. If  $\sum_{i=M}^{T-1} u_i > A$ , then  $\frac{1}{1 + \sum_{i=M}^{T-1} u_i / A} < \frac{1}{2}$  so that term 1 is smaller than term 2 in (10). Therefore, it is optimal to then choose  $\beta_T^* = 0$ . Note that this choice is possible because  $\sum_{i=M}^{T-1} u_i > A$  implies that  $M < T$ .

In this case, the value of (10) minimized over  $\beta_T$  equals  $\frac{1}{A + \sum_{i=M}^{T-1} u_i}$ .

2. If  $\sum_{i=M}^{T-1} u_i \leq A$ , then  $\frac{1}{1 + \sum_{i=M}^{T-1} u_i / A} \geq \frac{1}{2}$  so that an optimal choice is  $\beta_T^* \in (0, 1]$  with

$$\frac{1 - \beta_T^*}{1 - \beta_T^* + \sum_{i=M}^{T-1} u_i / A} \leq \frac{1}{2}. \tag{11}$$

Indeed, term 2 in (10) then equals  $1/2$  while term 1, which is equal to or greater than  $1/2$ , is not relevant because  $\beta_T^* > 0$ .

In this case, the value of (10) minimized over  $\beta_T$  equals  $\frac{1}{2A}$ .

Combining these two cases, we obtain

$$\begin{aligned}
& \min_{\beta_T} \frac{1}{A} \left( \frac{1}{1 + \sum_{i=M}^{T-1} u_i / A} \mathbb{1}_{\beta_T=0} + \mathbb{1}_{\beta_T \neq 0} \max \left\{ \frac{1 - \beta_T}{1 - \beta_T + \sum_{i=M}^{T-1} u_i / A}, \frac{1}{2} \right\} \right) \\
&= \frac{1}{A + \sum_{i=M}^{T-1} u_i} \mathbb{1}_{\sum_{i=M}^{T-1} u_i > A} + \frac{1}{2A} \mathbb{1}_{\sum_{i=M}^{T-1} u_i \leq A} \\
&= \frac{1}{A + \max\{\sum_{j=M}^{T-1} u_j, A\}}.
\end{aligned}$$

Consequently, (9) is reduced to (2). In the case of  $\beta_T^* = 0$ , we have  $M^* < T$ , and the optimal benchmark is VWAP with  $\beta_j^* = \frac{u_j}{\sum_{i=M}^{T-1} u_i}$ , as proven in Step 2a.

If  $\beta_T^* > 0$ , we saw above in the second case that the value of (10) minimized over  $\beta_T$

equals  $\frac{1}{2A}$  so that

$$\frac{c|V|}{A + \max\{\sum_{j=M}^{T-1} u_j, A\}} + Q(M) = \frac{c|V|}{2A} + Q(M)$$

on  $\beta_T^* > 0$ . Because  $Q$  is strictly decreasing, we obtain the optimal  $M^* = T$ , which then implies that  $\beta_T^* = 1$ . Note that condition (11) is satisfied for this choice. Therefore, the only optimal  $\beta_T^* > 0$  is  $\beta_T^* = 1$ , and we obtain the dichotomy presented in the main result: either all weight is in the auction or the benchmark is VWAP without an auction.

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